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# An Algorithm Based on Priority Rules for Solving a Multi-drone Routing Problem in Hazardous Waste Collection

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Abstract—This research investigates the problem of assigning pre-scheduled trips to multiple drones to collect hazardous waste from different sites in the minimum time. Each drone is subject to essential restrictions: maximum flying capacity and recharge operation. The goal is to assign the trips to the drones so that the waste is collected in the minimum time. This is done if the total flying time is equally distributed among the drones. An algorithm was developed to solve the problem. The algorithm is based on two main ideas: sort the trips according to a given priority rule and assign the current trip to the first available drone. Three different priority rules have been tested: Shortest Flying Time, Longest Flying Time, and Median Flying Time. Two recharging conditions are maintained: recharging needed time and recharging full duration. By applying each priority rule and each recharging condition, we generate a six versions of the algorithm. The six versions of the proposed algorithm were implemented in Java programming language. The results were analyzed and compared proving that the Longest Flying Time priority rule surpasses the other two rules. Moreover, recharging a drone just enough for taking the next trip proved to be better than fully recharging it.

Keywords—Drones; trip assignment; priority rules; flying capacity; load balance

## I. INTRODUCTION

The Vehicle Routing Problem (VRP) is an important problem discussing the process of finding the optimal routes for a single or multiple vehicles, to perform a certain task or service. VRPs have been extended for numerous of problems. Most recently, VRPs using a single or multiple drones also called Unmanned Aerial Vehicles (UAVs) have spread widely. Those varied to truck-drone, truck-multi-drone, and multi-truck-multi-drone problems to achieve precise goals and maintaining a set of constraints. The author in [1] presented the vehicle routing problem with drones (VRPD) as a design of the combination of truck-drone routes and timetables to serve a group of customers with specific requirements and time limitations. The applications of VRPD or as called, Drone Routing Problems (DRPs) received a massive amount of attention from different fields and concerned parties. Media has shown interest in drones equipped with cameras to cover stories in different areas [2], commerce companies have been intrigued with last-Mile delivery using UAVs for speeding the process and lowering shipping costs [3], scientists for great discoveries [4], medical staff for urgent transportation of equipment [5], program developers and even the entertainment sector for mesmerizing drones shows. Environmentalists have utilized drones for addressing great problems threatening our earth, affecting humans and all living creatures [6]. Specialized VRPs regarding critical environmental issues such as CO2e emission, air pollution and global warming are called Green Vehicle Routing Problems (G-VRPs). Hazardous waste is another serious problem that puts us in danger by the minute. Taking advantage of the VRPD and G-VRP, our study highlights the formation of a hazardous waste collection plan, where waste is collected in the minimum time possible using multiple drones (MTMD). The authors in [7] proved the strongly NP-hardness of the problem and proposed a 2phase approach and a linear program to solve it using a single Unmanned Aerial Vehicle. In their research, a set of efficient trips have been produced for each solved instance. Each trip is a path that starts from a depot site, visits some other sites to collect the maximum waste in the minimum time, respecting the drone flying and waste capacities before returning back to the depot for charging and some minor maintenance. In the current research, we adopt those datasets of generated trips and extend the study using multiple drones. The motivation behind this research was mainly to propose efficient algorithms that will allow an optimal or near-optimal use of the flying and weight capacities of a fleet of drones in order to collect hazardous waste in a shortest time from many sites. This research studies a new Drone Routing Problem (DRP) using mdrones. Constraints related to the flying capacity and recharge duration of each drone are set. The aim of the study is to assign pre-generated tours in [7] so that completion time of the last drone is minimum. To solve this problem, the authors proposed priority rules-based algorithm. These rules have the purpose of sorting the trips prior distribution among the drones. The algorithm allows a maximum use of each drone's charge without exceeding the flying capacity. An experimental study using benchmarks was conducted to compare the efficiency of the proposed algorithm for each priority rule.

The remaining of the paper is structured as follows. In Section II, the literature review highlights the related and relevant areas of the research. The methodology used for conducting the research is detailed in Section III. Next, the experimental results and discussion are explained in Section IV followed by the conclusion and future work.

## II. LITERATURE REVIEW

New Vehicle Routing Problems (VRPs) studies have been intensively emerging in the last few decades, especially those

using Unmanned Aerial Vehicles (UAVs). The author in [8] acknowledged that the interest in the use of drones in various applications has grown significantly in recent years. Accompanied with specific constraints, limitations and several goals, VRPs are becoming progressively challenging. Most commonly, Different approaches to solve Drone Routing Problems (DRPs) have been proposed for last-Mile delivery. Moreover, transportation of medical equipment as in [9], collection of waste, surveillance, search and rescue missions and even forest monitoring to detect fire as presented in [10]. The author in [11] stated that similar to the traditional vehicle routing problem, the vehicle routing problem with drones (VRPD) involves the use of drones in addition to trucks for delivery. When trucks and drones are routed together, the challenge is substantially more complex and distinct from traditional vehicle routing problems.

Most DRPs were organized to truck-drone routing problems, as [12] demonstrated how incorporating truck-drone tandems into transportation systems can enhance delivery speed while also allowing the fleet size to be reduced without impacting delivery times or adding to truck drivers' workload. The author in [13] presented and solved a truck-drone hybrid routing problem with time-dependent road travel time (TDHRP-TDRTT) to resolve the truck-drone cooperative issue. TDHRP-TDRTT was solved using an iterative local search heuristic method. Whereas, in [14], they thought of several drones, each with a different set of capabilities, such as speed and battery life, are transported by and sent out from the truck, working together to satisfy customers. The truck must wait until each drone returns and numerous drones can all be dispatched at once. The proposed model was given the name heterogeneous drone-truck routing problem (HDTRP), and a formulation of the problem using mixed-integer programming was provided.

Likewise, A parallel Drone Scheduling Traveling Salesman Problem was solved in [15]. In this approach, deliveries are divided between a truck and one or more drones. The truck goes on a tour from the depot, but the drones can only go back and forth. The goal is to finish as soon as possible. The problem was expanded by taking into account many vehicles. A hybrid metaheuristic, as well as a Mixed Integer Linear Programming formulation and a simple branch-and-cut algorithm, were also presented as solutions. In addition, A mixed-integer linear programming (MILP) model for the multi-trip Capacitated Arc Routing Problem was proposed in [16] to reduce the overall cost of waste collection (CARP).

Logically, the algorithms developed in solving the problems varied. A widely known algorithm is the neighborhood search. In [17], an Adaptive Large Neighborhood Search metaheuristic for the vehicle routing problem with drones (VRPD) was proposed, following minimum time and cost constraints, using multiple trucks. They developed a mathematical model to create a problem close to the Flying Sidekick Traveling Salesman Problem to optimize huge instances. Nevertheless, a tabu search heuristic algorithm in [18] proposed a new neighborhood generation process for a distribution company, to transport commodities from a single depot to multiple dealers. And a hybrid genetic algorithm was studied in [19] combining sweep and genetic algorithms in an enhanced approach. Both are recognized algorithms to solve the VRPD in the literature.

Other important algorithms are those whose called nature based. An ant colony optimization (ACO) technique was developed to solve the NP-hard VRPs with drones in [20]. The studies showed that the suggested ACO algorithm can solve the VRDP efficiently for diverse size instances and area distributions. Similarly, a novel dynamical artificial bee colony (DABC) was used in [21] to reduce operating costs. To identify employed bee swarm and onlooker bee swarm, two bee swarms were formed. Also, variable neighborhood descent was used in employed bee phase and onlooker bee phase in varying methods. An artificial bee colony-based hybrid approach was developed to solve the waste collecting problem while taking the halfway disposal pattern into account in [22]. Moreover, [23] suggested a Particle Swarm Optimization (PSO) algorithm to solve uncertain VRP, alongside a decoding scheme to improve its efficiency. According to [24], a simulated annealing heuristic algorithm was proposed to solve one of the Green Vehicle Routing Problems (G-VRPs), which is the Hybrid Vehicle Routing Problem (HVRP).

The clustering algorithms are majorly common. The author in [25] proposed two different hybrid metaheuristic algorithms in regards of the Clustered Vehicle Routing Problem (Clu-VRP). The first algorithm relies on an Iterated Local Search algorithm, in which only possible solutions are searched. The second one is a hybrid genetic search where the shortest Hamiltonian path between every set of vertices inside each cluster is precomputed.

The Location-Routing Problem for delivering orders to a group of clients was discussed in [26]. They attempted to reduce the overall CO<sub>2</sub> emissions using trucks and drones for last-mile deliveries. The problem was resolved using a mathematical model. Their findings focused on how adopting greener transportation technology can be a start to the substantial contribution that UAVs make to problems with parcel delivery routing. Hence, it is considered one of the G-VRPs, which highlight serious environmental issues like high CO<sub>2e</sub> emission, pollution, global warming, low sustainability lifestyle and many others, as clearly stated in [27]. Waste collection problem is another example of the G-VRPs. Conveniently, [28] introduced a particle swarm optimization (PSO) approach in a capacitated vehicle-routing problem (CVRP) model to establish the optimal waste collection and route optimization solutions. The PSO-based CVRP model included threshold waste level (TWL) and scheduling approaches. Solomon's insertion algorithm was developed to solve a real-world waste collection vehicle routing problem with time windows (VRPTW) in [29].

This research investigates the process of collecting hazardous waste from different sites using constrained drones (UAVs), proposing priority rules-based algorithm to help efficiently collect the waste in the minimum time using multiple drones (MTMD). Using the outcomes of [7] as the base of the study, extending it by assigning the trips to multiple drones instead of a single drone.

Despite the magnitude of research conducted on VRPs following versatile approaches and proposing applicable solutions, shedding lights on such problem with this significance and effect on the daily life and unknown future of livings on the face of earth is crucial and deserving. Also, employing the latest modern technology like drones to accomplish the desired goal efficiently makes a distinctive difference. On that account,

presenting an algorithm maintaining sensitive constraints and imposing priority rules as a solution would narrow the gap and expedite the plan to a safer, healthier and waste free planet.

To the best of our knowledge, there is no existing literature that addresses the specific challenges associated with using drones in hazardous waste collection. Therefore, we are unable to compare the proposed algorithm with existing solutions in this context.

### III. RESEARCH METHODOLOGY

To assign the trips to the drones in an efficient way so that the total waste will be collected from all sites in the minimum time, we developed an algorithm that is mainly based on three principles:

- Order the trips to be traveled according to their flying times. The trips are then executed according to either an increasing order of their flying times: Shortest Flying Time (SFT), a decreasing order of their flying times: Longest Flying Time (LFT), or alternatively from the median time: Median Flying Time (MFT).
- Assign the current trip to the first available drone.
  This principle will guarantee an efficient use of the
  drones, balance distributing the flying times between
  the drones, and collect the waste in a minimum time.
- When the current charge of a given drone is less than
  the flying duration of the next trip, then charge it
  according to two principles. The first is fully recharge
  the drone before taking the next trip. The second is to
  recharge the drone just enough for it to take the next
  trip.

The suggested algorithm enables multiple drones to access the list of trips at the same time, where each drone has a maximum flying and weight capacity, as well as a recharging duration. The flying capacity and drone's recharge are the two key constraints defined. The algorithm's two primary tasks are: parallel assignment of trips to multiple drones and recharging of drones. The trips are assigned to the drones as long as the list of trips is not empty, employing one of three proposed priority rules and following one of the charging principles. As a result, six versions of the main algorithm are proposed. The details of these algorithms are illustrated in Algorithms 1 to 4. Note that Algorithms 1 and 2 each has two versions. A version uses the SFT rule and another version uses the LFT rule. The notations used to develop the algorithm are summarized in Table I.

TABLE I. PROBLEM NOTATIONS

| n     | Number of trips in an instance  |
|-------|---|
| m     | Number of drones  |
| T     | Maximum flying time of a drone if fully charged                         |
| r     | Time needed to fully recharge a drone                                   |
| $c_j$ | Completion time of the last trip assigned to drone $j, 1 \leq j \leq m$ |
| $t_i$ | Flying duration of trip $i, 1 \leq i \leq n$                            |
| $q_j$ | Remaining flying capacity of drone j                                    |

Note that steps 3 to 5 in all algorithms have the purpose of initializing the completion times of the drones to zeros and

**Algorithm 1** Assignment of trips to drones using SFT/LFT rule and charging drones when needed

```
1: input: m, n, T, r, t_i
     output: C_{max}: The completion time of the last trip.
    Sort t_i in increasing/decreasing order
 3: for j \leftarrow 0 to m-1 do
      c_i \leftarrow 0
        q_j \leftarrow T
 6: end for
 7: i \leftarrow 0
 8: while (i < n) do
        a \leftarrow index \ of \min c_i
        if (q_a < t_i) then
10:
           c_a \leftarrow \frac{t_i - q_a}{T} \times r + t_i
q_a \leftarrow 0
12:
13:
14:
           c_a \leftarrow c_a + t_i
       q_a \leftarrow q_a - t_i end if
15:
16:
17:
       i \leftarrow i + 1
18: end while
19: return \max c_i
```

Algorithm 2 Assignment of trips to drones using SFT/LFT rule and charging drones fully

```
1: input: m, n, T, r, t_i
     output: C_{max}: The completion time of the last trip.
 2: Sort t_i in increasing/decreasing order
 3: for j \leftarrow 0 to m-1 do
 4:
        c_j \leftarrow 0
        q_j \leftarrow T
 5:
 6: end for
 7: i \leftarrow 0
 8: while (i < n) do
         a \leftarrow index \ of \ \min c_i
          \begin{array}{c} \textbf{if } (q_a < t_i) \textbf{ then} \\ c_a \leftarrow \frac{T - q_a}{T} \times r + t_i \end{array} 
10:
11:
12:
13:
            c_a \leftarrow c_a + t_i
14:
         end if
15:
         q_a \leftarrow q_a - t_i
         i \leftarrow i + 1
17: end while
18: return \max c_i
```

set their initial charge to full so that they can fly for T unites of time. Step 9 assigns the next trip to the first free drone. Step 10 in Algorithms 1 and 2 and 17 in Algorithms 3 and 4 means that the drone does not have enough charge to fly the next trip. For that reason, the drone needs either to be partially enough charged to fly the next trip (as in algorithms 1 and 3) or fully charged up to its maximum flying capacity T (as in algorithms 2 and 4). The last statement in all algorithms returns the maximum completion flying time of the last trip for all drones. The six proposed algorithms ensure an early completion time of the waste collection by assigning the next trip to the first available drone. This method maintains also an equilibrium between the total flying time between the drones.

# **Algorithm 3** Assignment of trips to drones using MFT rule and charging drones when needed

```
1: input: m, n, T, r, t_i
     output: C_{max}: The completion time of the last trip.
 2: Sort t_i in increasing order
 3: for j \leftarrow 0 to m-1 do
        c_i \leftarrow 0
         q_j \leftarrow T
 6: end for
 7: left \leftarrow \lfloor \frac{n}{2} \rfloor, right \leftarrow \lfloor \frac{n}{2} \rfloor + 1, k \leftarrow 0
 8: while (k < n) do
        a \leftarrow index \ of \ \min c_i
        if (k \mod 2 = 0) then
10:
            i \leftarrow left
11:
12:
            left \leftarrow left - 1
13:
14:
            i \leftarrow right
            right \leftarrow right + 1
15:
         end if
16:
        if (q_a < t_i) then
17:
            c_a \leftarrow \frac{t_i - q_a}{T} \times r + t_i
q_a \leftarrow 0
18:
19:
20:
21:
            c_a \leftarrow c_a + t_i
            q_a \leftarrow q_a - t_i
22:
         end if
23:
         k \leftarrow k + 1
24:
25: end while
26: return \max c_i
```

# **Algorithm 4** Assignment of trips to drones using MFT rule and and charging drones fully

```
1: input: m, n, T, r, t_i
     output: C_{max}: The completion time of the last trip.
 2: Sort t_i in increasing order
 3: for j \leftarrow 0 to m-1 do
        c_j \leftarrow 0 \\ q_j \leftarrow T
 5:
 6: end for
 7: left \leftarrow \lfloor \frac{n}{2} \rfloor, right \leftarrow \lfloor \frac{n}{2} \rfloor + 1, k \leftarrow 0
 8: while (k < n) do
 9:
         a \leftarrow index \ of \ \min c_i
10:
         if (k \mod 2 = 0) then
11:
            i \leftarrow left
            left \leftarrow left - 1
12:
         else
13:
             i \leftarrow right
14:
15:
            right \leftarrow right + 1
16:
         if (q_a < t_i) then c_a \leftarrow \frac{T-q_a}{T} \times r + t_i
17:
18:
19:
             c_a \leftarrow c_a + t_i
20:
21:
         end if
22:
         q_a \leftarrow q_a - t_i
         k \leftarrow k + 1
23:
24: end while
25: return \max c_i
```

#### IV. EXPERIMENTAL STUDY

The six versions of the proposed algorithms were implemented in Java. Distributing the number of trips among multiple drones grants completing all the trips in a sooner time than assigning all the trips to a single drone, because all the drones are flying simultaneously; whenever a drone completes a trip and it is available in the depot, the program assigns another trip to it, not taking into account the status of the other drones.

Testing the program using the three priority rules: SFT, LFT and MFT while recharging only for the time needed once, and again recharging the drones till they are full, generated six distinctive cases. Twenty instances varying from 101 to 941 sites, with a number of trips in the range [23, 249], have been used for validation of the proposed algorithms. A summary of results is shown in Tables II, III, IV, and V. For the six algorithms, the time each drone takes to complete the trips assigned to it was compared; the longest time means that it is the last drone to finish. As observed, the last drone finishing time of the cases that follow the LFT rule and recharge only for the time needed, is mostly the shortest, which means the LFT rule frequently guarantees a faster execution of the routing process. However, in some cases, multiple drones finish on the same time.

TABLE II. Values of  $C_{max}$  for Different Instances when m=2

|           |     | Ne   | eded cha | rge  | Full charge |      |      |  |
|-----------|-----|------|----------|------|-------------|------|------|--|
| Instance  | n   | SFT  | LFT      | MFT  | SFT         | LFT  | MFT  |  |
| $I_{101}$ | 23  | 294  | 289      | 284  | 296         | 293  | 284  |  |
| $I_{121}$ | 28  | 392  | 384      | 392  | 393         | 388  | 393  |  |
| $I_{141}$ | 32  | 416  | 409      | 416  | 416         | 411  | 416  |  |
| $I_{161}$ | 34  | 353  | 345      | 353  | 353         | 349  | 353  |  |
| $I_{181}$ | 52  | 606  | 598      | 606  | 607         | 600  | 607  |  |
| $I_{341}$ | 69  | 725  | 719      | 720  | 725         | 720  | 720  |  |
| $I_{301}$ | 75  | 708  | 702      | 699  | 708         | 703  | 703  |  |
| $I_{381}$ | 82  | 960  | 952      | 960  | 960         | 957  | 957  |  |
| $I_{321}$ | 87  | 1107 | 1101     | 1097 | 1076        | 1089 | 1098 |  |
| $I_{361}$ | 100 | 866  | 862      | 866  | 866         | 862  | 866  |  |
| $I_{461}$ | 118 | 1365 | 1357     | 1365 | 1365        | 1357 | 1365 |  |
| $I_{481}$ | 127 | 960  | 1128     | 1133 | 1365        | 1357 | 1133 |  |
| $I_{641}$ | 121 | 1494 | 1470     | 1488 | 1494        | 1471 | 1488 |  |
| $I_{661}$ | 143 | 1847 | 1840     | 1840 | 1847        | 1841 | 1841 |  |
| $I_{701}$ | 156 | 1519 | 1511     | 1519 | 1519        | 1512 | 1519 |  |
| $I_{721}$ | 162 | 1662 | 1652     | 1662 | 1662        | 1658 | 1662 |  |
| $I_{741}$ | 183 | 1493 | 1487     | 1489 | 1493        | 1488 | 1489 |  |
| $I_{861}$ | 194 | 2188 | 2182     | 2188 | 2188        | 2184 | 2188 |  |
| $I_{901}$ | 207 | 2158 | 2151     | 2153 | 2158        | 2152 | 2153 |  |
| $I_{941}$ | 249 | 2159 | 2151     | 2154 | 2159        | 2152 | 2154 |  |

Moreover, the three cases where the drones are fully recharged every time the battery ran out, took a slightly longer total time to complete all the trips, regardless of the number of drones used, than those three cases which the drones are recharged for only the time needed to be assigned the next suitable trip. Table VI confirms this result for instance  $I_{101}$ .

Although, the recharge count could be the same for both recharging conditions, the remaining charge in the cases that recharge only the needed duration is always zero. Whereas the cases that recharge fully, have useless remaining charge; that adds extra period of time, which increases the completion time of all the trips. Therefore, it is better to consider the cases that will reduce the total time by charging the drones only for the time needed.

TABLE III. Values of  $C_{max}$  for Different Instances when m=3

|           |     | Needed charge |      |      | Full charge |      |      |  |
|-----------|-----|---------------|------|------|-------------|------|------|--|
| Instance  | n   | SFT           | LFT  | MFT  | SFT         | LFT  | MFT  |  |
| $I_{101}$ | 23  | 196           | 191  | 195  | 198         | 195  | 197  |  |
| $I_{121}$ | 28  | 272           | 266  | 264  | 273         | 270  | 264  |  |
| $I_{141}$ | 32  | 284           | 277  | 283  | 284         | 279  | 284  |  |
| $I_{161}$ | 34  | 246           | 234  | 242  | 246         | 237  | 242  |  |
| $I_{181}$ | 52  | 419           | 403  | 408  | 420         | 408  | 410  |  |
| $I_{341}$ | 69  | 483           | 474  | 484  | 483         | 476  | 485  |  |
| $I_{301}$ | 75  | 467           | 462  | 471  | 470         | 465  | 469  |  |
| $I_{381}$ | 82  | 651           | 640  | 647  | 651         | 642  | 646  |  |
| $I_{321}$ | 87  | 729           | 726  | 735  | 730         | 729  | 736  |  |
| $I_{361}$ | 100 | 591           | 581  | 584  | 591         | 582  | 584  |  |
| $I_{461}$ | 118 | 923           | 911  | 921  | 923         | 911  | 918  |  |
| $I_{481}$ | 127 | 767           | 752  | 761  | 767         | 754  | 762  |  |
| $I_{641}$ | 121 | 1004          | 995  | 1000 | 1004        | 995  | 1000 |  |
| $I_{661}$ | 143 | 1227          | 1226 | 1230 | 1233        | 1229 | 1232 |  |
| $I_{701}$ | 156 | 1015          | 1007 | 1018 | 1015        | 1007 | 1020 |  |
| $I_{721}$ | 162 | 1112          | 1098 | 1114 | 1112        | 1102 | 1113 |  |
| $I_{741}$ | 183 | 997           | 987  | 1000 | 997         | 988  | 1000 |  |
| $I_{861}$ | 194 | 1465          | 1457 | 1466 | 1464        | 1458 | 1465 |  |
| $I_{901}$ | 207 | 1437          | 1430 | 1443 | 1438        | 1431 | 1442 |  |
| $I_{941}$ | 249 | 1444          | 1431 | 1445 | 1444        | 1431 | 1444 |  |

TABLE IV. Values of  $C_{max}$  for Different Instances when m=4

|           |     | Ne   | eded cha | rge  | Full charge |      |      |
|-----------|-----|------|----------|------|-------------|------|------|
| Instance  | n   | SFT  | LFT      | MFT  | SFT         | LFT  | MFT  |
| $I_{101}$ | 23  | 148  | 142      | 150  | 150         | 143  | 151  |
| $I_{121}$ | 28  | 200  | 190      | 201  | 200         | 193  | 201  |
| $I_{141}$ | 32  | 212  | 203      | 215  | 212         | 205  | 215  |
| $I_{161}$ | 34  | 191  | 173      | 186  | 191         | 181  | 185  |
| $I_{181}$ | 52  | 309  | 301      | 312  | 309         | 305  | 313  |
| $I_{341}$ | 69  | 373  | 363      | 368  | 373         | 366  | 368  |
| $I_{301}$ | 75  | 352  | 346      | 352  | 356         | 349  | 361  |
| $I_{381}$ | 82  | 488  | 480      | 486  | 488         | 482  | 487  |
| $I_{321}$ | 87  | 554  | 548      | 554  | 555         | 551  | 555  |
| $I_{361}$ | 100 | 440  | 431      | 444  | 441         | 432  | 444  |
| $I_{461}$ | 118 | 692  | 681      | 691  | 692         | 682  | 692  |
| $I_{481}$ | 127 | 578  | 562      | 577  | 578         | 564  | 577  |
| $I_{641}$ | 121 | 761  | 749      | 755  | 761         | 749  | 755  |
| $I_{661}$ | 143 | 925  | 919      | 928  | 927         | 920  | 928  |
| $I_{701}$ | 156 | 765  | 754      | 767  | 766         | 756  | 768  |
| $I_{721}$ | 162 | 838  | 827      | 841  | 838         | 832  | 842  |
| $I_{741}$ | 183 | 748  | 742      | 755  | 748         | 744  | 754  |
| $I_{861}$ | 194 | 1103 | 1095     | 1102 | 1104        | 1098 | 1102 |
| $I_{901}$ | 207 | 1085 | 1075     | 1085 | 1085        | 1075 | 1085 |
| $I_{941}$ | 249 | 1091 | 1079     | 1087 | 1091        | 1083 | 1088 |

TABLE V. Values of  $C_{max}$  for Different Instances when m=5

|           |     | Needed charge |     |     | Full charge |     |     |  |
|-----------|-----|---------------|-----|-----|-------------|-----|-----|--|
| Instance  | n   | SFT           | LFT | MFT | SFT         | LFT | MFT |  |
| $I_{101}$ | 23  | 125           | 114 | 125 | 126         | 118 | 126 |  |
| $I_{121}$ | 28  | 169           | 159 | 167 | 169         | 162 | 167 |  |
| $I_{141}$ | 32  | 181           | 170 | 176 | 181         | 172 | 176 |  |
| $I_{161}$ | 34  | 155           | 135 | 155 | 155         | 136 | 154 |  |
| $I_{181}$ | 52  | 261           | 242 | 257 | 262         | 247 | 258 |  |
| $I_{341}$ | 69  | 294           | 286 | 299 | 294         | 287 | 299 |  |
| $I_{301}$ | 75  | 281           | 274 | 287 | 284         | 278 | 292 |  |
| $I_{381}$ | 82  | 398           | 384 | 394 | 398         | 390 | 399 |  |
| $I_{321}$ | 87  | 451           | 444 | 450 | 452         | 447 | 452 |  |
| $I_{361}$ | 100 | 355           | 345 | 360 | 355         | 345 | 360 |  |
| $I_{461}$ | 118 | 556           | 543 | 556 | 557         | 546 | 557 |  |
| $I_{481}$ | 127 | 471           | 453 | 468 | 471         | 455 | 467 |  |
| $I_{641}$ | 121 | 613           | 602 | 609 | 613         | 602 | 609 |  |
| $I_{661}$ | 143 | 750           | 738 | 750 | 750         | 739 | 749 |  |
| $I_{701}$ | 156 | 620           | 610 | 619 | 620         | 613 | 619 |  |
| $I_{721}$ | 162 | 676           | 662 | 673 | 676         | 663 | 671 |  |
| $I_{741}$ | 183 | 603           | 595 | 605 | 603         | 596 | 604 |  |
| $I_{861}$ | 194 | 883           | 872 | 885 | 883         | 876 | 883 |  |
| $I_{901}$ | 207 | 876           | 865 | 872 | 876         | 866 | 872 |  |
| $I_{941}$ | 249 | 876           | 860 | 877 | 874         | 860 | 876 |  |

TABLE VI. TOTAL TIME COMPARISON FOR INSTANCE  $I_{101}$ 

|   | Ne          | eded cha | arge | Full charge |     |     |  |
|---|-------------|----------|------|-------------|-----|-----|--|
| m | SFT LFT MFT |          |      | SFT         | LFT | MFT |  |
| 2 | 562         | 563      | 562  | 566         | 569 | 568 |  |
| 3 | 556         | 557      | 556  | 561         | 566 | 559 |  |
| 4 | 549         | 550      | 550  | 557         | 557 | 556 |  |
| 5 | 544         | 543      | 544  | 553         | 558 | 553 |  |

In addition, the average of the time each drone took to complete the assigned trips was calculated. The difference between the time each drone took, and the average was also calculated, to illustrate how balanced the assigning of trips among the drones was. The total difference of those differences was then summed up to accentuate the exact gap between the time all the drones took to complete the trips and the average of it. It is noticeable that the total difference generated by SFT rule is always the biggest compared to the LFT and MFT rules. Whilst, the MFT cases often have smaller total difference than the SFT cases and a few of them have the same total difference as the LFT or SFT cases. In numerous instances, the MFT cases have the smallest total difference out of all three rules. Furthermore, a decreasing pattern in the number of those instances was noticeable, where the number of instances and the number of drones used have an inverse relationship; whenever the number of drones used increases, the number of the instances having the smallest total difference in the MFT cases decreases. However, the LFT cases mostly have the smallest total difference amongst all three. Additionally, in some LFT cases the total time of all the drones is the same as the average time, which means the total difference is zero. This comparison truly demonstrated how the LFT rule performs best in regards of the assignment of trips among the drones, despite their number, which ensures that all the drones take up nearly the same time to complete the trips. Moreover, it is an indication that the drones visit approximately the same number of trips. As a resultant, sorting the trips from the longest flying time to the shortest flying time (LFT) almost equally assigns the trips to the drones, which executes a more efficient program, effectively fulfilling the objective of the proposed algorithm MTMD.

"Fig. 1, "Fig. 2", "Fig. 3" and "Fig. 4" below, clarify the contrast between the total differences in each rule (SFT, LFT and MFT) for all 20 instances using 2, 3, 4 and 5 drones.

## V. CONCLUSION AND FUTURE WORK

#### A. Conclusion

In conclusion, studying the Vehicle Routing Problem using Drones (VRPD) and how it can be applied to help finding a solution to the accumulated hazardous waste problem was the first step. The diversified literature of VRPs gave the proper background and knowledge to investigate this major problem and propose an algorithm to resolve it.

Adopting the outcomes of an existent research as a dataset, six versions of a proposed algorithm were brought forward that allow multiple drones to access the list of trips at the same time, where the drones have maximum flying and weight capacities and a recharge duration. Concentrating on the flying capacity and the recharging of drones as constraints, the

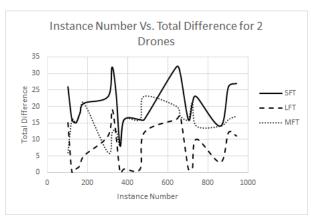


Fig. 1. Total differences for each rule using 2 drones.

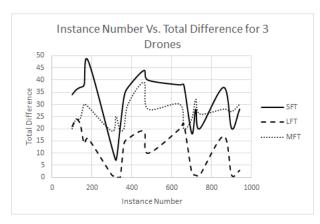


Fig. 2. Total differences for each rule using 3 drones.

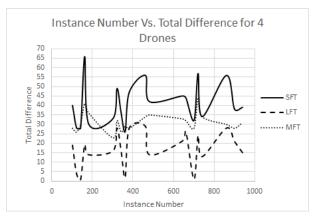


Fig. 3. Total differences for each rule using 4 drones.

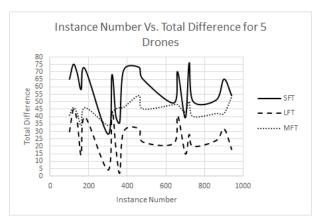


Fig. 4. Total differences for each rule using 5 drones.

presented algorithms handle the routing through two main tasks: the assignment of trips to multiple drones and recharging of drones. It suggested three priority rules in sorting the flying times before assigning them to the drones. The first rule is sorting from the shortest to the longest flying time (SFT) in ascending order. The second rule is sorting from the longest to the shortest flying time (LFT) in descending order. The third rule is sorting the flying times in ascending or descending order first and then starting from the median trip time (MFT). Following each rule, whenever the drones need recharging, two conditions were taken into consideration: recharging the drones just for the time needed to be assigned the next adequate trip or recharging the drones fully, taking their total recharge duration. Hence, the algorithms map routes in minimum time using multiple drones (MTMD).

Moreover, the algorithm's versions were implemented using Java to program was developed validating the algorithm, which produced six cases that test each individual rule (SFT, LFT and MFT) with both recharging conditions (needed and full).

After various testing of the six cases, the output of the program resulted in valuable findings. The LFT rule was proven the current best in assignment of trips to the multiple drones. It almost distributes the trips equally, which means approximately the same total flying time, recharge count and trip count of the drones. The program was tested using common several testing types as well, to assure its effectiveness in following the MTMD algorithm.

## B. Future Work

Aggregating the work done, positive improvements are expected to better the research, starting from developing a program with higher performance and dedicating enough time to test all the instances of the dataset using a greater number of drones.

In addition, the Drone Routing Problem (DRP) is progressing continuously; researchers are studying new challenging extensions of it momentarily. Consequently, proposing a better algorithm to develop an enhanced program as a solution to this substantial problem is an aspiration kept in mind. An existing similar problem called the Subset-Sum problem could

be beneficial for evolving the rules the MTMD algorithm proposed.

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